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Inline assembly for x86 in Linux

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Putting the pieces together

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Bharata B. Rao offers a guide to the overall use and structure of inline assembly for x86 on the Linux platform. He covers the basics of inline assembly and its various usages, gives some basic inline assembly coding guidelines, and explains the instances of inline assembly code in the Linux kernel.

If you're a Linux kernel developer, you probably find yourself coding highly architecture-dependent functions or optimizing a code path pretty often. And you probably do this by inserting assembly language instructions into the middle of C statements (a method otherwise known as inline assembly). Let's take a look at the specific usage of inline assembly in Linux. (We'll limit our discussion to the IA32 assembly.)

GNU assembler syntax in brief

Let's first look at the basic assembler syntax used in Linux. GCC, the GNU C Compiler for Linux, uses AT&T assembly syntax. Some of the basic rules of this syntax are listed below. (The list is by no means complete; I've included only those rules pertinent to inline assembly.)

Register naming

Register names are prefixed by %. That is, if eax has to be used, it should be used as %eax.

Source and destination ordering

In any instruction, source comes first and destination follows. This differs from Intel syntax, where source comes after destination.

mov %eax, %ebx, transfers the contents of eax to ebx.

Size of operand

The instructions are suffixed by b, w, or I, depending on whether the operand is a byte, word, or long. This is not mandatory; GCC tries provide the appropriate suffix by reading the operands. But specifying the suffixes manually improves the code readability and eliminates the possibility of the compilers guessing incorrectly.

movb %al, %bl -- Byte move
movw %ax, %bx -- Word move
movl %eax, %ebx -- Longword move

Immediate operand

An immediate operand is specified by using \$.

movl \$0xffff, %eax -- will move the value of 0xffff into eax register.

Indirect memory reference

Any indirect references to memory are done by using ().

movb (%esi), %al -- will transfer the byte in the memory pointed by esi into al register

Inline assembly

GCC provides the special construct "asm" for inline assembly, which has the following format:

GCC's "asm" construct

asm (assembler template	
: output operands	(optional)
: input operands	(optional)
: list of clobbered registers	(optional)
);	

In this example, the assembler template consists of assembly instructions. The input operands are the C expressions that serve as input operands to the instructions. The output operands are the C expressions on which the output of the assembly instructions will be performed.

Inline assembly is important primarily because of its ability to operate and make its output visible on C variables. Because of this capability, "asm" works as an interface between the assembly instructions and the C program that contains it.

One very basic but important distinction is that *simple inline assembly* consists only of instructions, whereas *extended inline assembly* consists of operands. To illustrate, consider the following example:

The basics of the inline assembly

```
{
    int a=10, b;
    asm ("movl %1, %%eax;
    movl %%eax, %0;"
        :"=r"(b) /* output */
        :"r"(a) /* input */
        :"%eax"); /* clobbered register */
}
```

In our example we have now made the value of "b" equal to "a" using assembly instructions. Note the following points:

- "b" is the output operand, referred to by %0 and "a" is the input operand, referred to by %1.
- "r" is a constraint on the operands, which specifies that variables "a" and "b" are to be stored in registers. Note that output operand constraint should have a constraint modifier "=" along with it to specify that it is an output operand.
- To use a register %eax inside "asm", %eax should be preceded by one more %, in other words, %%eax, as "asm" uses %0, %1 etc. to identify variables. Anything with single % will be treated as input/output operands and not as registers.

- The clobbered register %eax after the third colon tells GCC that the value of %eax is to be modified inside "asm", so GCC won't use this register to store any other value.
- movl %1, %%eax moves the value of "a" into %eax, and movl %%eax, %0 moves the contents of %eax into "b".
- When the execution of "asm" is complete, "b" will reflect the updated value, as it is specified as an output operand. In other words, the change made to "b" inside "asm" is supposed to be reflected outside the "asm".

Now let's go into each of these items in a bit more detail.

Assembler template

The assembler template is a set of assembly instructions (either a single instruction or a group of instructions) that gets inserted inside the C program. Either each instruction should be enclosed within double quotes, or the entire group of instructions should be within double quotes. Each instruction should also end with a delimiter. The valid delimiters are newline(\n) and semicolon(;). '\n' may be followed by a tab(\t) as a formatter to increase the readability of the instructions GCC generates in the assembly file. The instructions refer to the C expressions (which are specified as operands) by numbers %0, %1 etc.

If you want to make sure that the compiler does not optimize the instructions inside the "asm", you can use the "volatile" keyword after "asm". If the program has to be ANSI C compatible, __asm__ and __volatile__ should be used instead of asm and volatile.

Operands

C expressions serve as operands for the assembly instructions inside "asm". Operands are the main feature of inline assembly in cases where assembly instructions do a meaningful job by operating on the C expressions of a C program.

Each operand is specified by a string of operand constraints followed by the C expression in brackets, for example: "constraint" (C expression). The primary function of the operand constraint is to determine the addressing mode of the operand.

You may use more than one operand both in input and output sections. Each operand is separated by a comma.

The operands are referenced by numbers inside the assembler template. If there are a total of n operands (both input and output inclusive), then the first output operand is numbered 0, continuing in increasing order, and the last input operand is numbered n-1. The total number of operands is limited to ten, or to the maximum number of operands in any instruction pattern in the machine description, whichever is greater.

List of clobbered registers

If the instructions in "asm" refer to hardware registers, we can tell GCC that we will use and modify them ourselves. GCC will consequently not assume that the values it loads into these registers will be valid. It isn't generally necessary to list input and output registers as clobbered because GCC knows that "asm" uses them (because they are specified explicitly as constraints). However, if the instructions use any other registers, implicitly or explicitly (and the registers are not present either in input or in the output constraint list), then the registers have to be specified as a clobbered list. Clobbered registers are listed after the third colon by specifying the register names as strings.

As far as keywords are concerned, if the instructions modify the memory in some unpredictable fashion, and not explicitly, then the "memory" keyword may be added to the list of clobbered registers. This tells GCC not to keep the memory values cached in the register across the instructions.

Operand constraints

As mentioned before, each operand in "asm" should be described by a string of operand constraints followed by the C expression in brackets. Operand constraints essentially determine the addressing mode of the operand in the instructions. Constraints can also specify:

- Whether an operand is allowed to be in a register, and which kinds of registers it may be included in
- Whether the operand can be a memory reference, and which kinds of addresses to use in such a case
- Whether the operand may be an immediate constant

Constraints may also require two operands to match.

Commonly used constraints

Of the available operand constraints, only a few are used with any frequency; these are listed below along with brief descriptions. For a complete list of operand constraints, refer to the GCC and GAS manuals.

Register operand constraint(r)

When operands are specified using this constraint, they get stored in General Purpose Registers. Take the following example:

```
asm ("movl %%cr3, %0\n" :"=r"(cr3val));
```

Here the variable cr3val is kept in a register, the value of %cr3 is copied onto that register, and the value of cr3val is updated into the memory from this register. When the "r" constraint is specified, GCC may keep the variable cr3val in any of the available GPRs. To specify the register, you must directly specify the register names by using specific register constraints.

a %eax b %ebx c %ecx d %edx S %esi D %edi

Memory operand constraint(m)

When the operands are in the memory, any operations performed on them will occur directly in the memory location, as opposed to register constraints, which first store the value in a register to be modified and then write it back to the memory location. But register constraints are usually used only when they are absolutely necessary for an instruction or they significantly speed up the process. Memory constraints can be used most efficiently in cases where a C variable needs to be updated inside "asm" and you really don't want to use a register to hold its value. For example, the value of idtr is stored in the memory location loc:

```
("sidt %0\n" : :"m"(loc));
```

Matching(Digit) constraints

In some cases, a single variable may serve as both the input and the output operand. Such cases may be specified in "asm" by using matching constraints.

```
asm ("incl %0" :"=a"(var):"0"(var));
```

In our example for matching constraints, the register %eax is used as both the input and the output variable. var input is read to %eax and updated %eax is stored in var again after increment. "0" here specifies the same constraint as the 0th output variable. That is, it specifies that the output instance of var should be stored in %eax only. This constraint can be used:

- In cases where input is read from a variable or the variable is modified and modification is written back to the same variable
- In cases where separate instances of input and output operands are not necessary

The most important effect of using matching restraints is that they lead to the efficient use of available registers.

Examples of common inline assembly usage

The following examples illustrate usage through different operand constraints. There are too many constraints to give examples for each one, but these are the most frequently used constraint types.

"asm" and the register constraint "r"

Let's first take a look at "asm" with the register constraint 'r'. Our example shows how GCC allocates registers, and how it updates the value of output variables.

In this example, the value of x is copied to y inside "asm". x and y are passed to "asm" by being stored in registers. The assembly code generated for this example looks like this:

```
main:
        pushl %ebp
        movl %esp,%ebp
        subl $8,%esp
        movl $10,-4(%ebp)
        movl -4(%ebp),%edx
                               /* x=10 is stored in %edx */
        /* asm starts here */
#APP
        movl %edx, %eax
                                /* x is moved to %eax */
        movl %eax, %edx
                                /* y is allocated in edx and updated */
#NO APP /* asm ends here */
        movl %edx,-8(%ebp)
                                /* value of y in stack is updated with
                                   the value in %edx */
```

GCC is free here to allocate any register when the "r" constraint is used. In our example it chose %edx for storing x. After reading the value of x in %edx, it allocated the same register for y.

Since y is specified in the output operand section, the updated value in %edx is stored in -8(%ebp), the location of y on stack. If y were specified in the input section, the value of y on stack would not be updated, even though it does get updated in the temporary register storage of y(%edx).

And since %eax is specified in the clobbered list, GCC doesn't use it anywhere else to store data.

Both input x and output y were allocated in the same %edx register, assuming that inputs are consumed before outputs are produced. Note that if you have a lot of instructions, this may not be the case. To make sure that input and output are allocated in different registers, we can specify the & constraint modifier. Here is our example with the constraint modifier added.

And here is the assembly code generated for this example, from which it is evident that x and y have been stored in different registers across "asm".

```
main:
    pushl %ebp
    movl %esp,%ebp
    subl $8,%esp
    movl $10,-4(%ebp)
    movl -4(%ebp),%ecx /* x, the input is in %ecx */
#APP
    movl %ecx, %eax
    movl %eax, %edx /* y, the output is in %edx */
#NO_APP
```

movl %edx,-8(%ebp)

Use of specific register constraints

Now let's take a look at how to specify individual registers as constraints for the operands. In the following example, the cpuid instruction takes the input in the %eax register and gives output in four registers: %eax, %ebx, %ecx, %edx. The input to cpuid (the variable "op") is passed to "asm" in the eax register, as cpuid expects it to. The a, b, c, and d constraints are used in the output to collect the values in the four registers, respectively.

```
asm ("cpuid"

: "=a" (_eax),

"=b" (_ebx),

"=c" (_ecx),

"=d" (_edx)

: "a" (op));
```

And below you can see the generated assembly code for this (assuming the _eax, _ebx, etc.... variables are stored on stack):

The strcpy function can be implemented using the "S" and "D" constraints in the following manner:

The source pointer src is put into %esi by using the "S" constraint, and the destination pointer dst is put into %edi using the "D" constraint. The count value is put into %ecx as it is needed by rep prefix.

And here you can see another constraint that uses the two registers %eax and %edx to combine two 32-bit values and generate a 64-bit value:

```
#define rdtscll(val) \
    __asm____volatile__ ("rdtsc" : "=A" (val))
The generated assembly looks like this (if val has a 64 bit memory space).
#APP
    rdtsc
#NO_APP
    movl %eax,-8(%ebp) /* As a result of A constraint
    movl %edx,-4(%ebp) %eax and %edx serve as outputs */
Note here that the values in %edx:%eax serve as 64 bit output.
```

Using matching constraints

Here you can see the code for the system call, with four parameters:

In the above example, four arguments to the system call are put into %ebx, %ecx, %edx, and %esi by using the constraints b, c, d, and S. Note that the "=a" constraint is used in the output so that the return value of the system call, which is in %eax, is put into the variable ___res. By using the matching constraint "0" as the first operand constraint in the input section, the syscall number ___NR_##name is put into %eax and serves as the input to the system call. Thus %eax serves here as both input and output register. No separate registers are used for this purpose. Note also that the input (syscall number) is consumed (used) before the output (the return value of syscall) is produced.

Use of memory operand constraint

Consider the following atomic decrement operation:

__asm___volatile__("lock; decl %0" :"=m" (counter) :"m" (counter));

The generated assembly for this would look something like this:

```
#APP
    lock
    decl -24(%ebp) /* counter is modified on its memory location */
#NO_APP.
```

You might think of using the register constraint here for the counter. If you do, the value of the counter must first be copied on to a register, decremented, and then updated to its memory. But then you lose the whole purpose of locking and atomicity, which clearly shows the necessity of using the memory constraint.

Using clobbered registers

Consider an elementary implementation of memory copy.

http://www-106.ibm.com/developerworks/library/l-ia.html (8 of 9) [4/2/2001 1:46:55 PM]

While lodsl modifies %eax, the lodsl and stosl instructions use it implicitly. And the %ecx register explicitly loads the count. But GCC won't know this unless we inform it, which is exactly what we do by including %eax and %ecx in the clobbered register set. Unless this is done, GCC assumes that %eax and %ecx are free, and it may decide to use them for storing other data. Note here that %esi and %edi are used by "asm", and are not in the clobbered list. This is because it has been declared that "asm" will use them in the input operand list. The bottom line here is that if a register is used inside "asm" (implicitly or explicitly), and it is not present in either the input or output operand list, you must list it as a clobbered register.

Conclusion

On the whole, inline assembly is huge and provides a lot of features that we did not even touch on here. But with a basic grasp of the material in this article, you should be able to start coding inline assembly on your own.

Resources

- Refer to the Using and Porting the GNU Compiler Collection (GCC) manual.
- Refer to the GNU Assembler (GAS) manual.
- Check out Brennan's Guide to Inline Assembly.

About the author

Bharata B. Rao has a bachelor of Engineering in Electronics and Communication from Mysore University, India. He has been working for IBM Global Services, India since 1999. He is a member of the IBM Linux Technology Center, where he concentrates primarily on Linux RAS (Reliability, Availability, and Serviceability). Other areas of interest are operating system internals and processor architecture. He can be reached at <u>rbharata@in.ibm.com</u>.

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